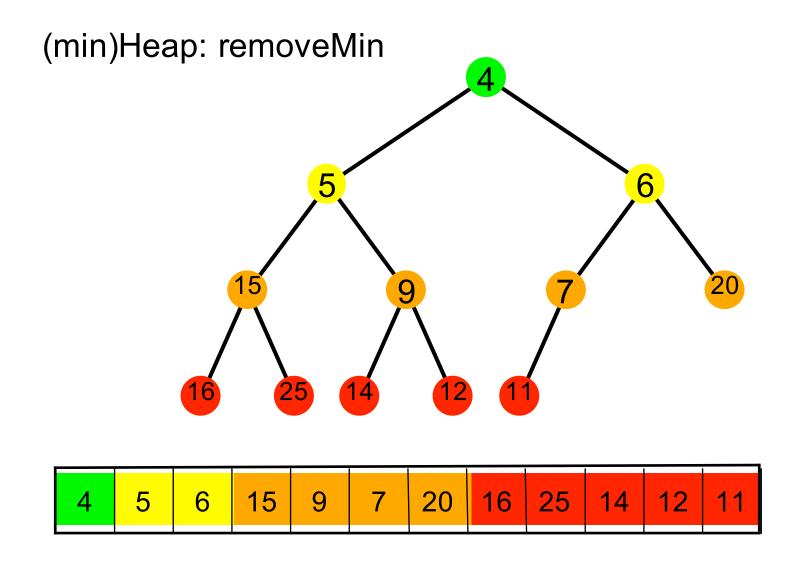
## Today's announcements:

MP6 available, due 11/17, 11:59p.

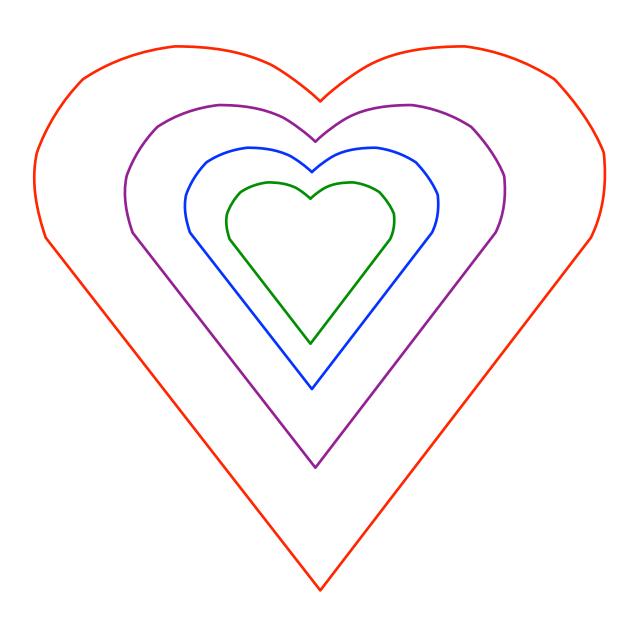


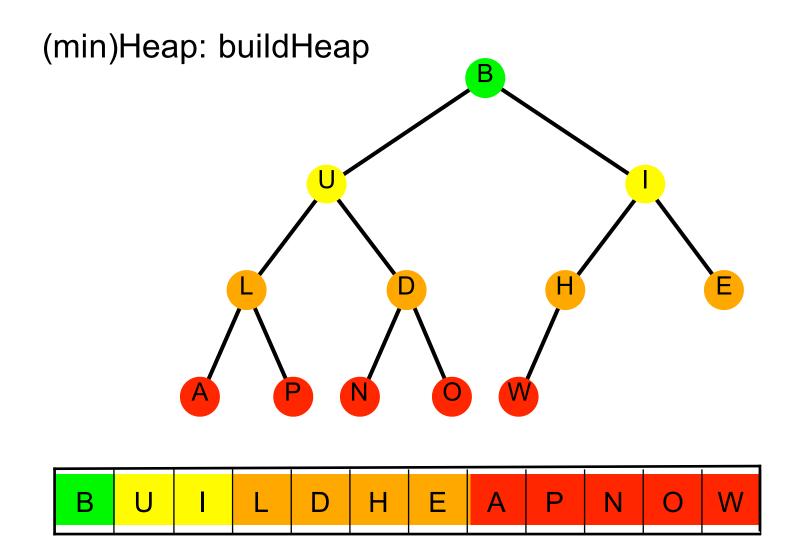
## Code:

```
template <class T>
T Heap<T>::removeMin() {
    T minVal = items[1];
    items[1] = items[size];
    size--;
    heapifyDown(1);
    return minVal;
}
```

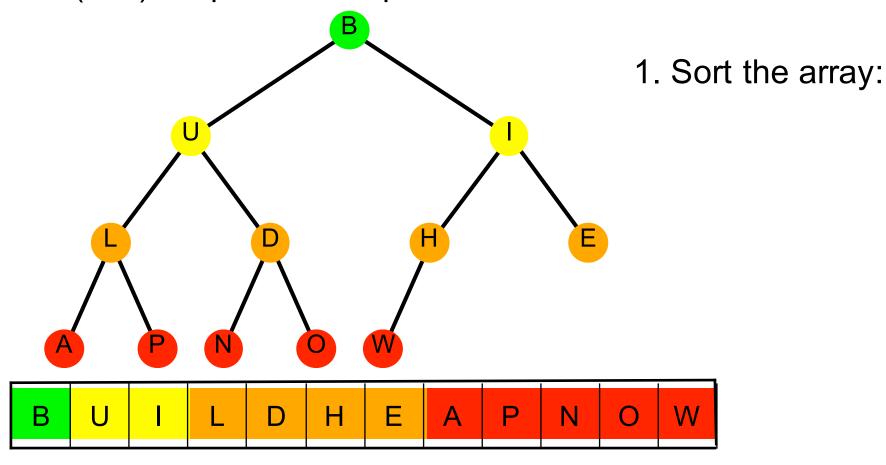
```
template <class T>
void Heap<T>::heapifyDown(int cIndex) {
   if (hasAChild(cIndex)) {
      minChildIndex = minChild(cIndex);
      if (items[cIndex] ____ items[minChildIndex] {
            swap(_____, ____);
            _____;
    }
}
```

## What have we done?



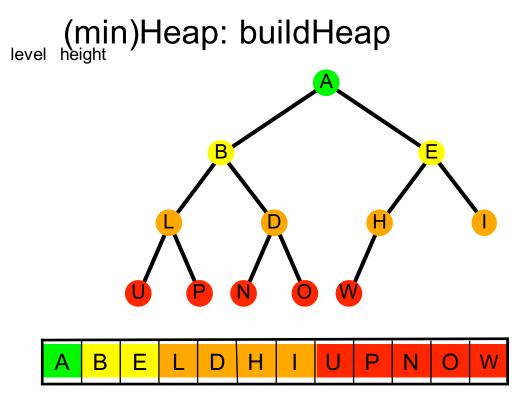


## (min)Heap: buildHeap - 3 alternatives



```
2. template <class T>
void Heap<T>::buildHeap() {
   for (int i=2;i<=size;i++)
       heapifyUp(i)
}</pre>
```

```
template <class T>
void Heap<T>::buildHeap() {
   for (int i=parent(size);i>0;i--)
       heapifyDown(i)
}
```



Thm: The running time of buildHeap on an array of size n is \_\_\_\_\_.

Instead of focussing specifically on running time, we observe that the time is proportional to the sum of the heights of all of the nodes, which we denote by S(h).

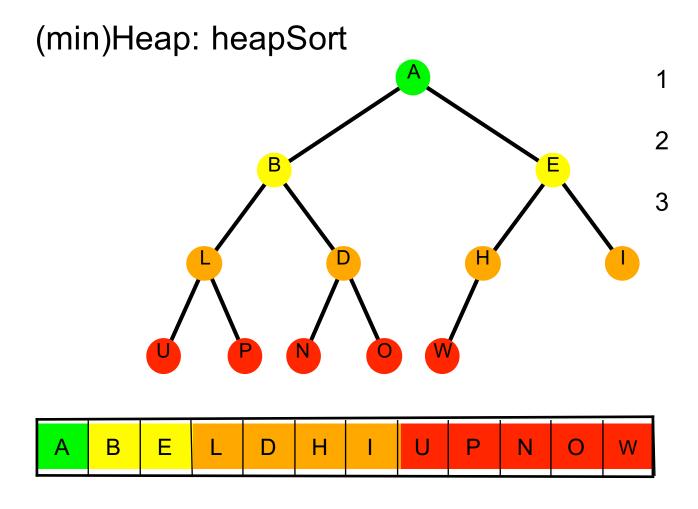
$$S(h) =$$

$$S(0) =$$

$$Soln S(h) =$$

Proof of solution to the recurrence:

But running times are reported in terms of n, the number of nodes...



Running time?

Why do we need another sorting algorithm?



This image reminds us of a	,	
which is one way we can implem	ent ADT	,
whose functions include	and	,
whose running times are		
This structure can be built in time	;	8
which helps us do a worst case t	ime sort	in place